

# A Hybrid Parallel Iterative Solver for Indefinite Systems in Interior Eigenvalue Computations

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Knowledge for Tomorrow

# Outline

Graphene simulation and the FEAST method

The CGMN algorithm

Parallelization

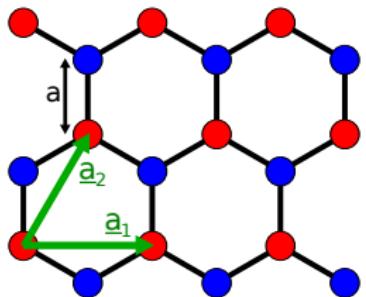
Experiments



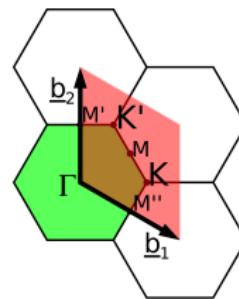
# Graphene simulation and the FEAST method



## Graphene



Physical space: carbon atoms in  
2D hexagonal mesh

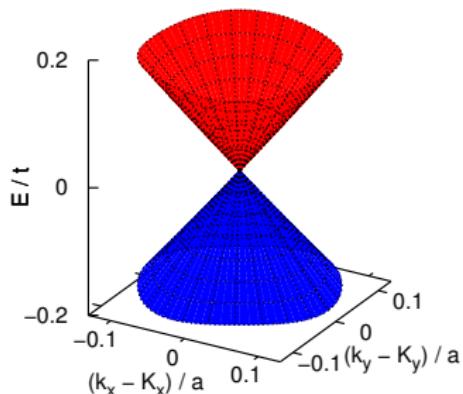
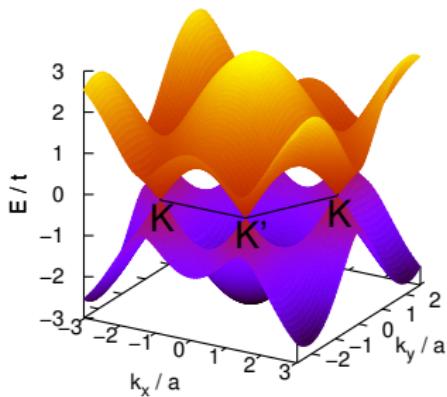


Fourier space ('reciprocal mesh')

Tight-binding Hamiltonian

$$H = -t \sum_{\langle ij \rangle} (c_i^\dagger c_j + c_j^\dagger c_i)$$

## Graphene (2)



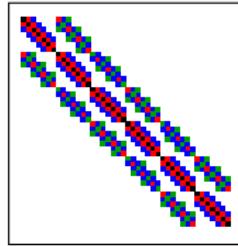
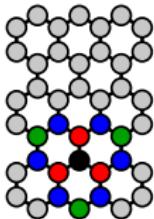
- Analytical solution for infinite Graphene sheet
- Dirac cones: graphene between conductor and semi-conducter

# Graphene modelling

- disorder
- long range stencil
- bilayer
- gate-defined quantum dots
- spin-orbit coupling
- ...

Long range Hamiltonian:

$$H = \sum_i V_i c_i^\dagger c_i - t \sum_{\langle ij \rangle} (c_i^\dagger c_j + c_j^\dagger c_i) - t' \sum_{\langle\langle ij \rangle\rangle} (c_i^\dagger c_j + c_j^\dagger c_i) - t'' \sum_{\langle\langle\langle ij \rangle\rangle\rangle} (c_i^\dagger c_j + c_j^\dagger c_i)$$



# The FEAST algorithm at a glance

**Input:**  $I_\lambda := [\underline{\lambda}, \bar{\lambda}]$ , an estimate  $\tilde{m}$  of the number of eigenvalues in  $I_\lambda$ .

**Output**  $\hat{m} \leq \tilde{m}$  eigenpairs with eigenvalue in  $I_\lambda$ .

**Perform:**

- ① Choose  $Y \in \mathbb{C}^{n \times \tilde{m}}$  of full rank and compute  

$$U := \frac{1}{2\pi i} \int_C (zB - A)^{-1} B dz Y,$$
- ② Form  $A_U := U^* A U$ ,  $B_U := U^* B U$ ,
- ③ Solve size- $\tilde{m}$  eigenproblem  $A_U \tilde{W} = B_U \tilde{W}^\sim$ ,
- ④ Compute  $(\tilde{X}, \tilde{X} := U \cdot \tilde{W})$ ,
- ⑤ If no convergence: go to Step 1 with  $Y := \tilde{X}$ .



## Linear systems for FEAST/graphene

Tough:

- very large ( $N = 10^8 - 10^{14}$ )
- complex symmetric and completely indefinite
- random numbers on and around the diagonal
- spectrum essentially continuous
- shifts get very close to the spectrum

But also nice in some ways:

- 2D mesh, very sparse ( $\sim 10$  entries/row)
- multiple RHS/shift (block methods, recycling, ...)

We need  $\mathcal{O}(100)$  Eigenpairs  $\implies$  very computationally heavy...



# The CGMN algorithm



## An ancient row projection method

- Björck and Elfving, 1979
- CG on the ‘minimum norm’ problem,  $AA^T x = b$
- preconditioned by SSOR
- efficient row-wise formulation
- extremely robust:  $A$  may be singular, non-square etc.
- row scaling alleviates issue of ‘squared condition number’



# Kernel operation: KACZ sweep

In CRS (rptr,val,col):

```

1: compute nrms=|| $a_{i,:}$ ||22
2: for (i=0; i<n; i++) do
   // compute  $a_{i,:}x - b_i$ 
3:   scal=-b[i]
4:   for (j=rptr[i]; j<rptr[i+1]; j++) do
5:     scal+=val[j]*x[col[j]]
6:   end for
7:   scal/=nrms[i]
   // update x
8:   for (j=rptr[i]; j<rptr[i+1]; j++) do
9:     x[col[j]]-=omega*scal*val[j]
10:  end for
11: end for
```

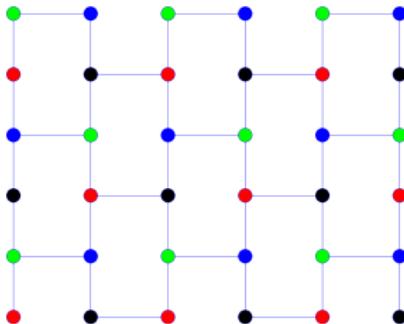
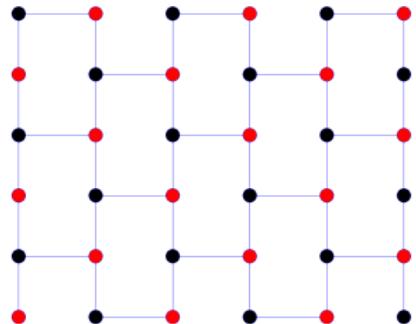
- Kaczmarz algorithm
- SOR( $\omega$ ) on the normal equations  $AA^T x = b$
- successive projections onto the hyperplanes defined by the rows of A



# Parallelization



## Multi-Coloring (MC)



- requires “distance 2” coloring
- software: ColPack  
<http://cscapes.cs.purdue.edu/coloringpage/software.htm>



## Component-Averaged Row Projection (CARP)

- Gordon & Gordon, 2005
- Kaczmarz locally
- write to halo
- exchange and average

Equivalent to Kaczmarz on a superspace of  $\mathbb{R}^n$



## Hybrid method: MC\_CARP-CG

- global MC would require...
  - an extremely scalable coloring method
  - very well-balanced colors
  - many global sync-points (> 20 colors in our examples)
- global CARP would require...
  - huge number of MPI procs
  - increasing amount of 'interior halo elements'
  - non-trivial implementation on GPU and Xeon Phi
  - increasing number of iterations

Idea: node-local MC with MPI-based CARP between the nodes



# Experiments



## Experimental setup

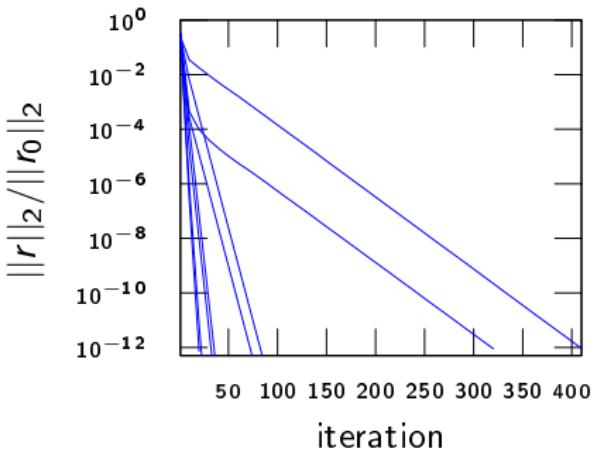
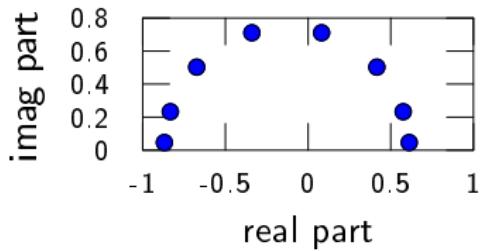
- Machine: Intel Xeon “Ivy Bridge”
- 10 cores/socket, 2 sockets/node
- InfiniBand between nodes

Here's what we do:

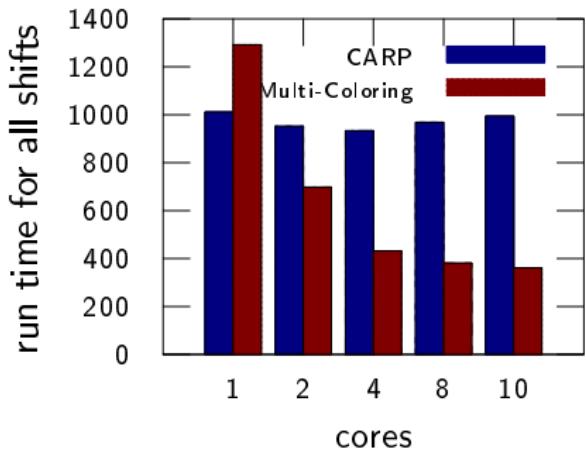
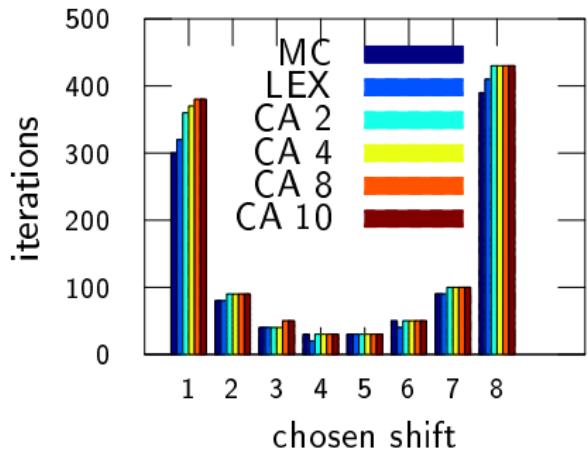
- pick some shifts that may occur in FEAST
- handle 8 rhs at once (for good performance)
- conv tol  $10^{-12}$
- solve linear systems using CGMN variants



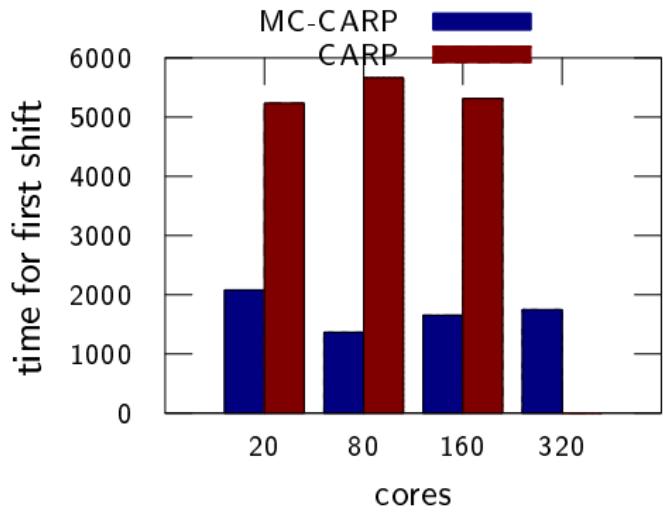
## Sequential CGMN for various shifts



## Coloring vs. CARP: single socket ( $1024^2$ dof)



## Weak scaling of Hybrid vs. CARP ( $4096^2$ dof/node)



## The (almost) final slide

- Graphene gives nice and challenging test cases for Lin. Alg.
- FEAST requires fast linear solvers for indef. systems
- row projection methods work very well here
- hybrid is a natural choice here - and works

Future work:

- integration in FEAST loop
- stencil-based implementation
- GPU and Xeon Phi

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## References

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